

From finite-system entropy to entropy rate for a Hidden Markov Process

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Abstract

A recent result presented the expansion for the entropy rate of a Hidden Markov Process (HMP) as a power series in the noise variable ϵ . The coefficients of the expansion around the noiseless ($\epsilon = 0$) limit were calculated up to 11th order, using a conjecture that relates the entropy rate of a HMP to the entropy of a process of finite length (which is calculated analytically). In this paper we generalize and prove the conjecture, and discuss its theoretical and practical consequences.

1 Introduction

Let $\{X_N\}$ be a finite state stationary Markov process over the alphabet $\Sigma = \{1, \dots, s\}$, and let $\{Y_N\}$ be its noisy observation (on the same alphabet). The process Y is generated by the Markov transition matrix $M = M_{s \times s} = \{m_{ij}\}$ and the emission matrix $I + \epsilon T$, where I is the $s \times s$ identity matrix, the matrix $T = T_{s \times s} = \{t_{ij}\}$ satisfies $t_{ii} < 0$, $t_{ij} \geq 0$, $\forall i \neq j$ and $\sum_{j=1}^s t_{ij} = 0$, and $\epsilon > 0$ is some constant. (There is no loss of generality here, as any stochastic matrix can be represented as $I + \epsilon T$.) This yields the probabilities $P(X_{N+1} = j | X_N = i) = m_{ij}$ and $P(Y_N = j | X_N = i) = \delta_{ij} + \epsilon t_{ij}$, where δ is Kronecker's delta. We consider the case of high signal to noise ratio ('High-SNR'), characterized by small values of ϵ , and assume strictly positive M ($m_{ij} > 0$) with a unique stationary distribution.

The process Y can be viewed as an observation of X through a noisy channel. It is a *Hidden Markov Process (HMP)*, governed by the parameters M , T and ϵ . *HMPs* have a rich

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theory, with applications in various fields, such as speech recognition ([1]), information theory ([2]) and signal processing ([3]). While we concentrate on a finite-state first-order *HMP*, our results can be easily generalized to more cases (e.g. continuous observations).

An important quantity for a stochastic process is the Shannon entropy rate, which measures its 'uncertainty per-symbol' ([4]). More formally, for $i \leq j$ let $[Y]_i^j$ denote the vector (Y_i, \dots, Y_j) . The entropy rate of Y is defined as:

$$\bar{H}(Y) = \lim_{N \rightarrow \infty} \frac{H([Y]_1^N)}{N} \quad (1)$$

Where $H(Y) = -\sum_Y P(Y) \log P(Y)$; Sometimes we omit the realization y of the variable Y , so $P(Y)$ should be understood as $P(Y = y)$. For a finite-entropy stationary process the limit (1) exists and \bar{H} can also be computed via the conditional entropy ([5]) as: $\bar{H}(Y) = \lim_{N \rightarrow \infty} H(Y_N | [Y]_1^{N-1})$. Here $H(U|V)$ represents the conditional entropy, which for random variables U and V is the average uncertainty of the conditional distribution of U given V , that is $H(U|V) = \sum_v P(V = v) H(U|V = v)$. By the entropy chain rule, it is also given as a difference of entropies, $H(U|V) = H(U, V) - H(V)$. This relation will be used below.

There is at present no explicit expression for the entropy rate of a *HMP* ([2, 6]). Few recent works ([6, 7, 8]) have studied the asymptotic behavior of \bar{H} in several regimes, albeit giving rigorously only bounds or at most second ([8]) order behavior. Here we generalize and prove a relationship, first posed in [8] as a conjecture, thereby turning the computation presented there, of \bar{H} as a series expansion up to 11th order in ϵ , into a rigorous statement.

2 Theorem Statement and Proof

We first state our main result, which will be proven at the end of the section.

Theorem 1 *Let $H_N \equiv H_N(M, T, \epsilon) = H([Y]_1^N)$ be the entropy of a system of length N , and let $C_N = H_N - H_{N-1}$. Let $B_\rho(0) \subset \mathbb{C}$ be some (complex) neighborhood of zero, in which the functions $\{C_N\}$ and \bar{H} are analytic in ϵ , with Taylor expansions given by:*

$$C_N(M, T, \epsilon) = \sum_{k=0}^{\infty} C_N^{(k)} \epsilon^k, \quad \bar{H}(M, T, \epsilon) = \sum_{k=0}^{\infty} C^{(k)} \epsilon^k \quad (2)$$

(The coefficients $C_N^{(k)}$ are functions of M and T . From now on we omit this dependence).

Then:

$$N \geq \left\lceil \frac{k+3}{2} \right\rceil \Rightarrow C_N^{(k)} = C^{(k)} \quad (3)$$

Analyticity of $\{C_N\}$ and \bar{H} around $\epsilon = 0$ was recently shown in [9]. One may also use [10], which showed that the law of the process Y is Gibbsian, together with the complete analyticity results for Gibbsian measures of [12], to deduce analyticity of \bar{H} . C_N is in fact an upper-bound ([5]) for \bar{H} . The behavior stated in Thm. 1 was discovered using symbolic computations, but was proven only for $k \leq 2$, in the binary symmetric case ([8]). Although it may appear technically involved, our proof is based on two simple ideas.

First, we distinguish between the noise parameters at different sites. We thus consider a more general process $\{Z_N\}$, where Z_i 's emission matrix is $I + \epsilon_i T$. The process $\{Z_N\}$ is determined by M, T and $[\epsilon]_1^N$. We define the following functions:

$$F_N(M, T, [\epsilon]_1^N) = H([Z]_1^N) - H([Z]_1^{N-1}) \quad (4)$$

Setting all the ϵ_i 's equal reduces this to the Y process, and in particular $F_N(M, T, (\epsilon, \dots, \epsilon)) = C_N(\epsilon)$.

Second, we observe that if a particular ϵ_i is set to zero, we must have $Z_i = X_i$. Thus, conditioning back to the past is 'blocked'. This is used to prove:

Lemma 1 *If $\epsilon_j = 0$ for some $1 < j < N$, then:*

$$F_N([\epsilon]_1^N) = F_{N-j+1}([\epsilon]_j^N) \quad (5)$$

Proof F can be written as the sum:

$$F_N = - \sum_{[Z]_1^N} \left[P([Z]_1^{N-1}) P(Z_N | [Z]_1^{N-1}) \log P(Z_N | [Z]_1^{N-1}) \right] \quad (6)$$

The dependence on $[\epsilon]_1^N$ and M, T is hidden in the probabilities $P(\cdot)$. Since $\epsilon_j = 0$, we have

$X_j = Z_j$, and conditioning further to the past is 'blocked':

$$\epsilon_j = 0 \Rightarrow P(Z_N|[Z]_1^{N-1}) = P(Z_N|[Z]_j^{N-1}) \quad (7)$$

Substituting in eq. 6 gives:

$$\begin{aligned} F_N &= - \sum_{[Z]_1^N} \left[P([Z]_1^{N-1}) P(Z_N|[Z]_j^{N-1}) \log P(Z_N|[Z]_j^{N-1}) \right] = \\ &= - \sum_{[Z]_j^N} P([Z]_j^N) \log P(Z_N|[Z]_j^{N-1}) = F_{N-j+1} \end{aligned} \quad (8)$$

■

Let $\vec{k} = [k]_1^N$ be a vector with $k_i \in \{\mathbb{N} \cup 0\}$. Define its 'weight' as $\omega(\vec{k}) = \sum_{i=1}^N k_i$. Define also:

$$F_N^{\vec{k}} \equiv \left. \frac{\partial^{\omega(\vec{k})} F_N}{\partial \epsilon_1^{k_1}, \dots, \partial \epsilon_N^{k_N}} \right|_{\vec{\epsilon}=0} \quad (9)$$

$C_N^{(k)}$ is obtained by summing the contributions $F_N^{\vec{k}}$ of all the vectors \vec{k} 's with weight k :

$$C_N^{(k)} = \frac{1}{k!} \sum_{\vec{k}, \omega(\vec{k})=k} F_N^{\vec{k}} \quad (10)$$

The next lemma shows that many such \vec{k} 's give zero contribution to the sum:

Lemma 2 *Let $\vec{k} = [k]_1^N$. If $\exists i, j, 1 \leq i < j < N$, with $k_i \geq 1, k_j \leq 1$, then $F_N^{\vec{k}} = 0$.*

Proof *Assume first $k_j = 0$. Using lemma 1 we get:*

$$\begin{aligned} F_N^{\vec{k}} &\equiv \left. \frac{\partial^{\omega(\vec{k})} F_N([\epsilon]_1^N)}{\partial \epsilon_1^{k_1}, \dots, \partial \epsilon_N^{k_N}} \right|_{\vec{\epsilon}=0} = \left. \frac{\partial^{\omega(\vec{k})} F_{N-j+1}([\epsilon]_j^N)}{\partial \epsilon_1^{k_1}, \dots, \partial \epsilon_N^{k_N}} \right|_{\vec{\epsilon}=0} = \\ &= \left. \frac{\partial^{\omega(\vec{k})-1}}{\partial \epsilon_1^{k_1}, \dots, \partial \epsilon_i^{k_i-1}, \dots, \partial \epsilon_N^{k_N}} \left[\frac{\partial F_{N-j+1}([\epsilon]_j^N)}{\partial \epsilon_i} \right] \right|_{\vec{\epsilon}=0} = 0 \end{aligned} \quad (11)$$

Assume now $k_j = 1$. Write the probability of Z :

$$P([Z]_1^N) = \sum_{[X]_1^N} P([X]_1^N) P([Z]_1^N|[X]_1^N) = \sum_{[X]_1^N} P([X]_1^N) \prod_{i=1}^N (\delta_{X_i Z_i} + \epsilon_i t_{X_i Z_i}) \quad (12)$$

Let $[Z]_1^{N(j \rightarrow a)}$ denote the vector we get from $[Z]_1^N$ by changing Z_j to a (while keeping other coordinates). Differentiating with respect to ϵ_j gives (see [11] for more details):

$$\begin{aligned} \left. \frac{\partial P([Z]_1^N)}{\partial \epsilon_j} \right|_{\epsilon_j=0} &= \sum_{[X]_1^N} \left[P([X]_1^N) t_{X_j Z_j} \prod_{i \neq j} (\delta_{X_i Z_i} + \epsilon_i t_{X_i Z_i}) \right] \Big|_{\epsilon_j=0} = \\ & \left\{ \sum_{a=1}^s t_{a Z_j} P([Z]_1^{N(j \rightarrow a)}) \right\} \Big|_{\epsilon_j=0} \end{aligned} \quad (13)$$

By Bayes' rule $P(Z_N | [Z]_1^{N-1}) = \frac{P([Z]_1^N)}{P([Z]_1^{N-1})}$, we get:

$$\begin{aligned} \left. \frac{\partial P(Z_N | [Z]_1^{N-1})}{\partial \epsilon_j} \right|_{\epsilon_j=0} &= \frac{1}{P([Z]_1^{N-1})} \sum_{a=1}^s t_{a Z_j} \left[P([Z]_1^{N(j \rightarrow a)}) - \right. \\ & \left. P(Z_N | [Z]_1^{N-1}) P([Z]_1^{N-1(j \rightarrow a)}) \right] \Big|_{\epsilon_j=0} \end{aligned} \quad (14)$$

This gives:

$$\begin{aligned} \left. \frac{\partial [P([Z]_1^N) \log P(Z_N | [Z]_1^{N-1})]}{\partial \epsilon_j} \right|_{\epsilon_j=0} &= \sum_{a=1}^s t_{a Z_j} \left\{ P([Z]_1^{N(j \rightarrow a)}) \log P(Z_N | [Z]_1^{N-1}) + \right. \\ & \left. P([Z]_1^{N(j \rightarrow a)}) - P(Z_N | [Z]_1^{N-1}) P([Z]_1^{N-1(j \rightarrow a)}) \right\} \Big|_{\epsilon_j=0} \end{aligned} \quad (15)$$

And therefore:

$$\begin{aligned} \left. \frac{\partial F_N}{\partial \epsilon_j} \right|_{\epsilon_j=0} &= \\ & - \sum_{a=1}^s t_{a Z_j} \left\{ \sum_{[Z]_1^N} \left[P([Z]_1^{N(j \rightarrow a)}) \log P(Z_N | [Z]_1^{N-1}) - P(Z_N | [Z]_1^{N-1}) P([Z]_1^{N-1(j \rightarrow a)}) \right] \right\} \Big|_{\epsilon_j=0} = \\ & \left\{ - \sum_{a=1}^s t_{a Z_j} \sum_{[Z]_j^N} \left[P([Z]_j^{N(1 \rightarrow a)}) \log P(Z_N | [Z]_j^{N-1}) - P(Z_N | [Z]_j^{N-1}) P([Z]_j^{N-1(1 \rightarrow a)}) \right] \right\} \Big|_{\epsilon_1=0} \end{aligned} \quad (16)$$

The latter equality comes from using eq. (7), which 'blocks' the dependence backwards. Eq. (16) shows that ϵ_i does not appear in $\left. \frac{\partial F_N}{\partial \epsilon_j} \right|_{\epsilon_j=0}$ for $i < j$, therefore $\left. \frac{\partial^{k_i+1} F_N}{\partial \epsilon_i^{k_i} \partial \epsilon_j} \right|_{\epsilon_j=0} = 0$ and $F_N^{\vec{k}} = 0$. ■

Before proving Thm. 1, we show here that adding zeros to the left of \vec{k} leaves $F_N^{\vec{k}}$ unchanged:

Lemma 3 Let $\vec{k} = [k]_1^N$ with $k_1 \leq 1$. Denote $\vec{k}^{(r)}$ the concatenation of \vec{k} and r zeros to the left: $\vec{k}^{(r)} = \underbrace{(0, \dots, 0)}_r, k_1, \dots, k_N$. Then:

$$F_N^{\vec{k}} = F_{r+N}^{\vec{k}^{(r)}}, \forall r \in \mathbb{N} \quad (17)$$

Proof Assume first $k_1 = 0$. Using lemma 1 we get:

$$F_{r+N}^{\vec{k}^{(r)}}([\epsilon]_1^{r+N}) = \frac{\partial^{\omega(\vec{k}^{(r)})} F_{r+N}([\epsilon]_1^{r+N})}{\partial \epsilon_{r+2}^{k_2}, \dots, \partial \epsilon_{r+N}^{k_N}} \Big|_{\vec{\epsilon}=0} = \frac{\partial^{\omega(\vec{k})} F_N([\epsilon]_{r+1}^{r+N})}{\partial \epsilon_{r+2}^{k_2}, \dots, \partial \epsilon_{r+N}^{k_N}} \Big|_{\vec{\epsilon}=0} = F_N^{\vec{k}}([\epsilon]_{r+1}^{r+N}) \quad (18)$$

The case $k_1 = 1$ is reduced back to the case $k_1 = 0$ by taking the derivative. Using eqs. (16,18), we get:

$$\begin{aligned} F_{N+1}^{\vec{k}^{(1)}}([\epsilon]_1^{N+1}) &= \frac{\partial^{\omega(\vec{k})-1}}{\partial \epsilon_3^{k_2} \dots \partial \epsilon_{N+1}^{k_N}} \left[\frac{\partial F_{N+1}}{\partial \epsilon_2} \Big|_{\epsilon_2=0} \right] \Big|_{\vec{\epsilon}=0} = \\ &= \frac{\partial^{\omega(\vec{k})-1}}{\partial \epsilon_3^{k_2} \dots \partial \epsilon_{N+1}^{k_N}} \left\{ - \sum_{a=1}^s t_{aZ_2} \sum_{[Z]_1^{N+1}} \left[P([Z]_1^{N+1(2 \rightarrow a)}) \log P(Z_{N+1} | [Z]_1^N) - \right. \right. \\ &\quad \left. \left. P(Z_{N+1} | [Z]_1^N) P([Z]_1^{N(2 \rightarrow a)}) \right] \Big|_{\epsilon_2=0} \right\} \Big|_{[\epsilon]_1^{N+1}=0} = \\ &= \frac{\partial^{\omega(\vec{k})-1}}{\partial \epsilon_2^{k_2} \dots \partial \epsilon_N^{k_N}} \left\{ - \sum_{a=1}^s t_{aZ_2} \sum_{[Z]_1^N} \left[P([Z]_1^{N(1 \rightarrow a)}) \log P(Z_N | [Z]_1^{N-1}) - \right. \right. \\ &\quad \left. \left. P(Z_N | [Z]_1^{N-1}) P([Z]_1^{N(1 \rightarrow a)}) \right] \Big|_{\epsilon_1=0} \right\} \Big|_{[\epsilon]_1^N=0} = F_N^{\vec{k}}([\epsilon]_1^N) \end{aligned} \quad (19)$$

This proved the claim for $r = 1$. The claim for larger r 's follows by induction. \blacksquare

We are now ready to prove our main theorem, which follows directly from lemmas 2 and 3:

Proof (Thm. 1) Let $\vec{k} = [k]_1^N$ with $\omega(\vec{k}) = k$. Define its 'length' as $l(\vec{k}) = N+1 - \min_{k_i > 1} \{i\}$. It easily follows from lemma 2 that $F_N^{\vec{k}} \neq 0 \Rightarrow l(\vec{k}) \leq \lceil \frac{k+3}{2} \rceil - 1$. Thus, according to lemma 3, we have:

$$F_N^{\vec{k}} = F_{\lceil \frac{k+3}{2} \rceil}^{(k_{N-\lceil \frac{k+3}{2} \rceil+1}, \dots, k_N)} \quad (20)$$

for all \vec{k} 's in the sum. Summing over all $F_N^{\vec{k}}$ with the same 'weight' gives $C_N^{(k)} = C_{\lceil \frac{k+3}{2} \rceil}^{(k)}$, $\forall N > \lceil \frac{k+3}{2} \rceil$. But from the analyticity of C_N and \bar{H} near $\epsilon = 0$ it can be shown by induction that $\lim_{N \rightarrow \infty} C_N^{(k)} = C^{(k)}$, therefore $C_N^{(k)} = C^{(k)}$, $\forall N \geq \lceil \frac{k+3}{2} \rceil$. ■

3 Conclusion

The theorem proven above sheds light on the connection between finite and infinite chains, and gives a practical and straightforward way to compute the entropy rate as a series expansion in ϵ up to an arbitrary power. The surprising 'settling' of the expansion coefficients $C_N^{(k)} = C^{(k)}$ for $N \geq \lceil \frac{k+3}{2} \rceil$, holds for the entropy. For other functions involving only conditional probabilities (e.g. relative entropy between two *HMPs*) a weaker result holds: the coefficients 'settle' for $N \geq k + 2$. One can expand the entropy rate in several parameter regimes. As it turns out, exactly the same 'settling' as was proven in Thm. 1 happens in the 'almost memoryless' regime, where the transition matrix M is close to a matrix which makes the X_i 's i.i.d, (i.e. a matrix whose rows are identical). This and other regimes, as well as the analytic behavior of the *HMP* ([9]), will be discussed elsewhere.

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