# Pair-List with Universal Quantifiers

## Goal for today

- To discuss a problem for our approach to questions with *wh* phrases: we have argued that pair-list readings of such questions (W-pair list) results from a family of questions denotation. Our first goal is to challenge this proposal by observing that the arguments for such a denotation extend to cases where the higher *wh*-phrase is replaced by a universal quantifier (to ∀-pair-list). Although various authors have suggested a family of questions denotation for ∀-pair-list, the syntactic and semantics mechanisms we've postulated do not yield this result.
- To present a method for extending the proposal (following Pafel 1999, and Preuss 2001)

## 1. The Problem

The two sentences in (1) seem equivalent on the pair list reading (exhaustivity, point-wise uniqueness).

(1) a. Which girl read which book?b. Which book did every girl read?

Moreover, our arguments for families of questions seem to extend:

### (2) **Plural Agreement:**

Imagine that at the end of the school year (11-12th grade) the teacher meets with every student to discuss plans for the future.

- a. \*The questions she will ask, (namely) who has plans to apply to college, are critical for the advice she will give.
- b. ? The questions she will ask, (namely) which student will apply to which university, are critical for the advice she will give.
- b. ? The questions she will ask, (namely) to which university will every student apply, are critical for the advice she will give.

## (3) **Quantificational variability:**

- a. \*For the most part I would like to know who will vote for John in the upcoming elections.
- b. For the most part I would like to know who will vote for whom in the upcoming elections.
- c. For the most part I would like to know for whom every one of my friends will vote in the upcoming elections.

### [(4) **Exceptives:**

- a. \*I would like to know which one of my friends will vote for Scott Brown except for my neighbor Fred.
- b. I would like to know which one of my friends will vote for whom except for my neighbor Fred.
- c. I would like to know for whom every one of my friends will vote except for my neighbor Fred.]

The family of questions analysis was automatic in the case of multiple *wh*-phrases (W-pair-list). But it is far from so, in the case of  $\forall$ -pair-list

**My goal**: to present a family of questions analysis for  $\forall$ -pair-list (a version of a proposal made in Pafel 1999, and Preuss 2001).

## 2. Enrichment of Karttunen's mechanisms

## 2.1. Reminder

- (5) a.  $[[C_{int}]] = \lambda p_{\alpha} \cdot \lambda q_{\alpha} \cdot p = q$  (\*i.e., the relation of identity\*) b.  $[[which boy]] = [[some boy]] = \lambda P_{et} \cdot \exists x [x \text{ is a boy and } P(x) = 1]$
- (6) Which boy came? LF:  $\lambda p \text{ [which boy } \lambda x \text{ [[}C_{int} p\text{]} \lambda w. x came_w \text{]]}$ Denotation (in a world w<sup>0</sup>):  $\lambda p. \text{ [[some boy]]}^{w0} (\lambda x. [\lambda w. x came in w]=p)$
- (7) Which girl read which book? LF<sub>1</sub> (single occurrence of C<sub>int</sub> → simple question → unique answer): λp [which girl λx which book λy [[C<sub>int</sub> p] λw. x read<sub>w</sub> y]] Denotation (in a world w<sup>0</sup>): λp<sub>st</sub>. [[some girl]]<sup>w0</sup> (λx. [[some book]]<sup>w0</sup> (λy. p = λw. x read<sub>w</sub> y)) = λp<sub>st</sub>. ∃x∈[[girl]]<sup>w0</sup> ∃y∈[[book]]<sup>w0</sup>, s.t. p = λw. x read y in w

In set notation:  $\{\lambda w. x \text{ read } y \text{ in } w: y \in [[book]]^{w0} \& x \in [[girl]]^{w0}\}$ 

(8) Which girl read which book? LF<sub>2</sub> (involves two occurrences of C<sub>int</sub>  $\rightarrow$  multiple questions  $\rightarrow$  multiple answers):  $\lambda Q$  [which girl  $\lambda x$  [C<sub>int</sub> Q]  $\lambda p$  [which book  $\lambda y$  [[C<sub>int</sub> p]  $\lambda w. x \text{ read}_w y$ ]] Denotation (in a world w<sup>0</sup>):  $\lambda Q_{st,t}$ . [[some girl]]<sup>w0</sup> ( $\lambda x. Q = \lambda p_{st}$  [[some book]]<sup>w0</sup> ( $\lambda y. p = \lambda w. x \text{ read } y \text{ in } w$ )) =  $\lambda Q_{st,t}. \exists x \in [[girl]]^{w0} \text{ s.t.}$   $Q = \lambda p_{st}. \exists y \in [[book]]^{w0}, \text{ s.t.}$   $p = \lambda w. x \text{ read } y \text{ in } w$ In set notation: { $\{\lambda w. x \text{ read } y \text{ in } w: y \in [[book]]^{w0}$ }:  $x \in [[girl]]^{w0}$ }

## 2.2. Up to a higher type and back down with MIN

Is there a way for *which book every girl read?* to have the denotation in (8)?

Yes: with the addition of two extra-pieces to the structure:

-Null operators that can be merged in various positions and move.

-A covert minimization operator, MIN (Pafel)

(9) Which book did every girl read? LF<sub>1</sub> (single occurrence of  $C_{int}$  and no extra-machinery  $\rightarrow$  simple question  $\rightarrow$  unique answer):  $\lambda p$  [which book  $\lambda y$  [[C<sub>int</sub> p]  $\lambda w$ . every girl  $\lambda x$  x read<sub>w</sub> y]] Denotation (in a world  $w^0$ ):  $\lambda p_{st}$ . [[some book]]<sup>w0</sup> ( $\lambda y$ .  $p = \lambda w$ . every girl<sub>w/w0</sub> read<sub>w</sub> y)) =  $\lambda p_{st}$ .  $\exists x \in [[gir1]]^{w0}$ , s.t.  $p = \lambda w$ . every girl {in w, in w<sup>0</sup>} read y in w In set notation: { $\lambda$ w. every girl read y in w: y  $\in [book]^{w_0}$ } Which book did every girl read? (10)LF<sub>2</sub> (single occurrences of  $C_{int}$  + null operator movement + QR above C +MIN  $\rightarrow$ family of questions  $\rightarrow$  multiple answers): Min( $\lambda K_{\leq 0, t}$  [every girl  $\lambda x K \lambda p$  [which book  $\lambda y$  [[C<sub>int</sub> p]  $\lambda w. x read_w y$ ]] (where  $O = \langle st, t \rangle$ ) Denotation (in a world  $w^0$ ):  $[[Min]](\lambda K, [[every girl]]^{w0}(\lambda x, K(\lambda p_{st} [[some book]]^{w0}(\lambda y, p = \lambda w, x read y in w))$  $\llbracket Min \rrbracket (\lambda K. \llbracket every girl \rrbracket^{w0} (\lambda x. \{\lambda w. x read y in w: y \in \llbracket book \rrbracket^{w0}\} \in K)$  $[[Min]] (\{K: \forall x \in [[girl]]^{w_0} [\{\lambda w. x read y in w: y \in [[book]]^{w_0}\} \in K\}])$ the minimal set of questions that for every girl g has a member the question which book did g read?

 $= \{ \{ \lambda w. x read y in w: y \in [[book]]^{w0} \} : x \in [[girl]]^{w0} \}$ 

 $[[Min]](K_{<\alpha,t,\succ}) = \text{the } Q \in K, \text{ s.t. } \forall Q' \in K Q \subseteq Q' \text{ (undefin. if a unique Q doesn't exist)}$ 

### 3. The distribution of pair list and Szabolcsi's Observation

As Pafel points out, we can continue to account for the restriction on pair-list readings. With any quantifier other than a universal, *min* will not be defined.

### 4. If (only) wh-phrases were universal quantifiers

Then we would have an identical account for W- and  $\forall$ -pair-list

# More on Quantificational Variability

#### **Goal for today**

- To discuss a problem for my reliance on Lahiri in the context of an account based on Dayal. One of my arguments for families of questions was based on a very partial discussion of Lahiri's assumptions about QV. But can these assumptions live peacefully with Dayal, or at least with a Max<sub>inf</sub> pressuppositon?
- To show that things are simple for veridical predicates, but less so for non-veridical predicates.
- To work towards an account of QV with non-veridical predicates that would still rely on Dayal's presuppositions for questions.
- To bring up some of the problems for a Lahiri type account of QV, raised by Beck and Sharvit.

### 1. Veridical Predicates

#### 1.1. For veridical predicates it simple enough to modify Lahiri so that Ans is involved.

(11) For the most part he knows who came. For most  $p \in H$  s.t. p is entailed by  $Ans(H)(w^0)$  and Aomic(p,H), he knows p Where  $H = \{\lambda w. x \text{ came in } w: x \in [[person]]^{w_0}\}$  $Atomic(p,Q_{st,t}) \Leftrightarrow_{def} p \in Q \& \forall p' \in Q[p \subseteq p' \rightarrow p'=p]$  (p is a weakest member in Q)

### 1.2. An analogy with definite descriptions

(12)	a.	For the most part John knows who came. True in w iff for most p such that p is an atomic (weakest) true proposition of the form <i>x came</i> John knows in w that p.
	b.	1
		True in w iff
		For most x such that x is an atomic part of <i>the books</i> ,
		John read x in w.
(13)	a.	For the most part John knows who came. True in w iff
		for most p such that p is an atomic part of Ans([[who came]](w),
		John knows in w that p.
	b.	For the most part John read the books.
		True in w iff

For most x such that x is an atomic part of [[the books]](w),

John read x in w.

#### **1.3.** Analogous syntax and semantics<sup>1</sup>

- (14) a. [[For the most part]] $(p_{st})(Q_{st,t}) = 1$  iff for most q such that q is entailed by p and is an atomic (weakest) element in Domain(Q), Q(p) =1.
  - b. [[For the most part]](x<sub>e</sub>)(Q<sub>e,t</sub>) = 1 iff for most y, such that y is a part of x and is an atomic element in Domain(Q), Q(p) =1.
    (\*identical entries if *entailed by* is the relevant *part of* relation on D<sub>st</sub>\*)
- (12)' Logical Forms for (12) (QR Sportiche like + Trace Conversion)
  - a. For the most part the True who came

 $\lambda p$ . John knows the  $\lambda q$  [ $p \in \overline{\text{True}}$  who came & q = p]

b. For the most part the books  $\lambda x$ . John read the  $\lambda y$ . [y $\in$  books & y = x]

(16) a. [[The True 
$$Q_{st}$$
]] =  $\lambda w$ .the unique proposition  $p \in [[True Q]](w)$   
such that  $\forall q \in [[True Q]](w) \rightarrow \{w: p \in [[True Q]](w)\} \subseteq \{w: q \in [[True Q]](w)\}$   
=  $\lambda w$ . the unique  $p \in [[Q]]$  such that  $p(w)=1$  and  $\forall q \in [[Q]](q(w)$   
=  $1 \rightarrow p \subseteq q$   
=  $\lambda w$ . Max<sub>inf</sub>(Q)(w)

- (12)" Logical Forms for (12)a (QR Sportiche like + Trace Conversion + world variables)  $\lambda w$ . For the most part [the True who came](w)  $\lambda p$ . he knows<sub>w</sub> [the<sub>p</sub> True who came ](w)
- (17)  $[[\lambda p. John knows_w [the_p True who came ](w)]] = \lambda p: p(w) = 1 & p \in [[who came]]. In w John knows that p$

### 2. The problem of non-Veridical Predicates

 (18) a. For the most part John and Mary agree on who came.
 NOT True in w iff

<sup>&</sup>lt;sup>1</sup> Loosely based on Magri's (2007) claim that distributivity is based on a "sortal" part whole relation: <u>https://docs.google.com/viewer?a=v&pid=sites&srcid=ZGVmYXVsdGRvbWFpbnxtYWdyaWdyZ3xneDo2ODM</u> <u>3Yjk5Mjl1NTgwMDYz&pli=1</u>

for most p such that p is an atomic part of Ans([[who came]](w), John and Mary agree with each other in w that p.

 b. For the most part John is certain about who came. NOT True in w iff for most p such that p is an atomic part of Ans([[who came]](w), John is certain in w that p.

## 3. My Eventual Goal

To keep the analogy with definite descriptions and to connect the problem of non-veridical predicates to the problem we see in the following. In these examples we need the world argument of *the unicorns* to be bound locally despite quantificational variability, just like the word variable of *Ans/The-True* 

(19) While John was dozing he imagined that there were 50 unicorns in his garden. For the most part he wanted to ride the unicorns.

### 4. An alternative for now

- (20) Movement + Reconstruction: For the most part, he is certain who came. For the most part  $\lambda p$ . John is certain the<sub>p</sub> Filter who came.
- (21) Association with Presupposition: [[For the most part]]  $(Q_{st,t}) = 1$  iff for most p such that Atomic(p,Domain(Q)): [Q(p)=1]
- (22) Denotation of nuclear scope:
   [[λp. John is certain the<sub>p</sub> who came]]<sup>w0</sup> =
   λp: p∈H & John believes that p is possible & John believes that Filter(H) is defined. John is certain that p
   Where H={λw. x came<sub>w</sub>: x∈[[person]]<sup>w0</sup>}

### 5. A paraphrase (building on Egré-and-Spector)

### 5.1. First pass

We can't employ E&S directly: E&S involves existential quantification over worlds which yields a strong enough result by virtue of reliance on Ans-Strong. But this (rather ingenious) maneuver is no longer effective once QV is introduced.

(23) a. For the most part John and Mary agree on  $[_Q$  who came]. NOT

True in w iff **J**w' for most p such that p is an atomic member of **[**Q**]** entailed by

Ans-Strong([[Q]])(w), John and Mary agree with each other in w that p.

b. For the most part John is certain about who came. NOT

True in w iff  $\exists w'$  for most p such that p is an atomic member of [Q] entailed by Ans-Strong([Q])(w), John is certain in w that p.

## 5.2. An alternative to E&S

- John V Q is true in w0 iff
   ∀w J V Ans-weak(Q)(w) is defined → John V Ans-weak(Q)(w) is true in w0
- John sometimes, mostly, always, never V Q is true in w0 iff
   ∀w J V Ans-weak(Q)(w) is defined →
   [For some/most/all/no p an atomic member of {p∈Q: Ans(Q)(w) entails p}
   John V p is true in w0]

### 6. Goal – Not there yet

To derive this paraphrase from reasonable assumptions in syntax and semantics